

## Serie 6

**Exercise 1** *The Morse structure for storing a square sparse matrix of order  $n$ , with  $m$  nonzero entries, has three variants. In each case, the nonzero entries are stored by rows in an array of real numbers `double am[m]` and the column index of `am[k]` in the integer array `int ic[m]`, in the variable `ic[k]`. For finding the row index of `am[k]`, one defines an integer array `ir`, according to the three possibilities:*

1. `int ir[m]`, `ir[k]=row index of am[k]`,
2. `int ir[n]`, `ir[i]=ki`, `ki` being the location in `am` of the first nonzero entry of the  $i$ -th row,
3. `int ir[n]`, `ir[i]=ni`, `ni` being the number of nonzero entries of the  $i$ -th row.

For example, for the matrix:

$$A = \begin{pmatrix} a_{11} & 0 & 0 & 0 \\ 0 & a_{22} & a_{23} & 0 \\ 0 & a_{32} & 0 & a_{34} \end{pmatrix},$$

the three possibilities are:

| <code>am</code>     | <code>a<sub>11</sub></code> | <code>a<sub>22</sub></code> | <code>a<sub>23</sub></code> | <code>a<sub>32</sub></code> | <code>a<sub>34</sub></code> |
|---------------------|-----------------------------|-----------------------------|-----------------------------|-----------------------------|-----------------------------|
| <code>ic</code>     | <code>1</code>              | <code>2</code>              | <code>3</code>              | <code>2</code>              | <code>4</code>              |
| <code>ir (1)</code> | <code>1</code>              | <code>2</code>              | <code>2</code>              | <code>3</code>              | <code>3</code>              |
| <code>ir (2)</code> | <code>1</code>              | <code>2</code>              | <code>4</code>              | <code>-</code>              | <code>-</code>              |
| <code>ir (3)</code> | <code>1</code>              | <code>2</code>              | <code>2</code>              | <code>-</code>              | <code>-</code>              |

Write down (on paper) the C code doing the multiplication of this matrix by a vector `double x[n]`. To your opinion, what is the best variant and why ?

**Exercise 2** *We consider a square banded matrix  $A$  with a symmetric structure (but not necessarily symmetric), which has  $1 + 2k_b$  nonzero bands (included the main diagonal), but which can be far from the main diagonal (example: the matrix resulting from the discretization of the Laplacian of Exercise 2, Serie 4).*

1. *Set up a structure for storing only the nonzero bands of  $A$  and which allows to do the multiplication  $\mathbf{y} = \mathbf{A}\mathbf{x}$  "bandwise" (the inner loop runs on the entries of a band). One can be inspired by the structure of the Blas function `dgblmv`, in which the bands are stored by rows in a bidimensional array, the number of rows being equal to the number of bands.*
2. *Use your structure for the matrix of the Laplacian. Complete the program in the file `gblmv.c` for doing this multiplication and compare (CPU time) to the result obtained with `dgblmv`. Choose  $N = 80$  and 2000 iterations. Explain the difference.*